

QoS for the IP Multicast Backbone

A look at the enabling technologies -
ATM, MPLS, DiffServ.....



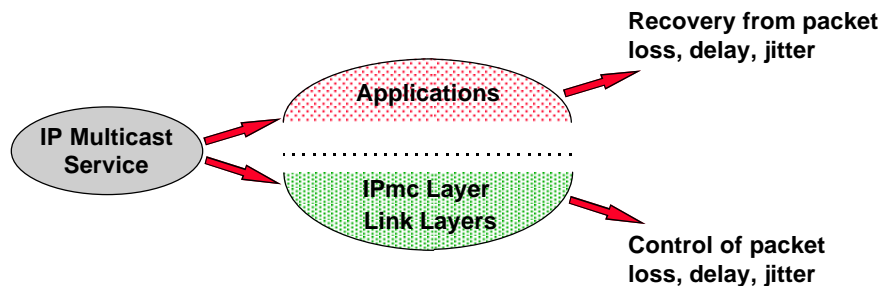
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What type of QoS are we talking about?

- “Quality” of service has two components
 - Above the IP layer
 - IP layer and below..... **we will focus here**



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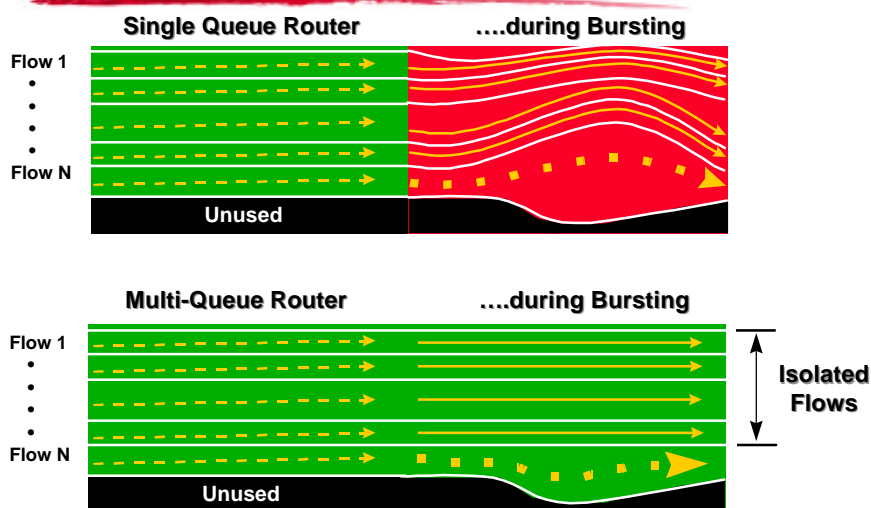
Backbone packet transport issue

- IP Packet streams vulnerable to congestion
 - Uncontrolled queuing and scheduling behavior during congestion leads to
loss delay jitter
- IPmc transport is often “2nd-class citizen”
 - Relegated to slower, ‘older’ routers
 - Tunneled through unicast networks
 - No protection from transient congestion caused by bursty unicast (revenue generating) traffic
- What will change this?
 - Line-rate isolation of unicast traffic and multicast traffic on common switching/routing technologies
 - Operator comfort level

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Isolating Unicast and Multicast traffic



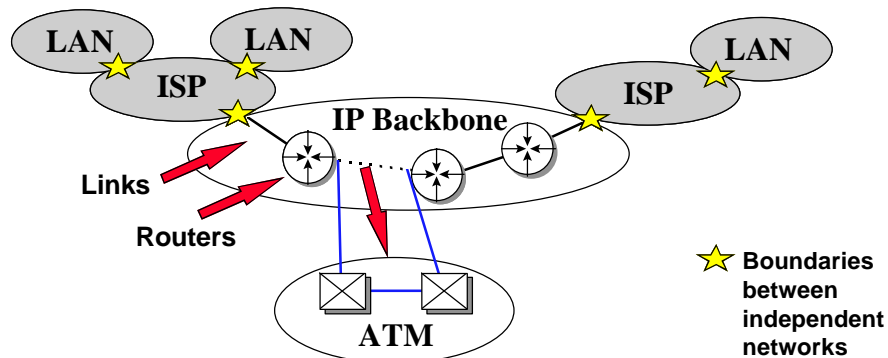
Multi-queue Routers Provide Isolation Between Traffic Classes

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Hierarchy of Network QoS

- End to End QoS depends on
 - Edge to Edge QoS
 - Router capabilities (e.g. diffserv, per-flow WFQ,....)
 - Link capabilities (e.g. ATM, MPLS/POS,...)



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Managing congestion

- Bursty and real-time traffic mix at routers/switches
 - Per-Router buffering absorbs transient overloads
 - Small buffers result in excessive packet loss
 - Large buffering leads to excessive delays
- *Need separate QUEUES (buffers) for traffic with different packet loss tolerance and burstiness*
- *Need traffic engineering to force packets along least congested paths*

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"CQS" - Classify, Queue, Schedule

Classify

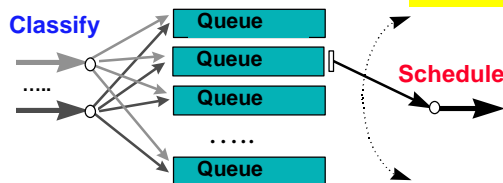
Establish context of the cells/packets

Queue

Absorb bursts, isolate customers

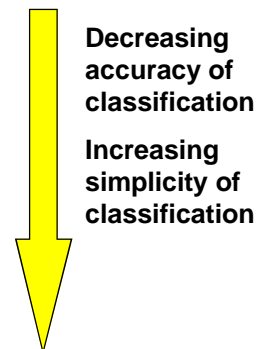
Schedule

Mediate access to link bandwidth
Ensure bounded delay/jitter



Establishing context

- A number of popular solutions differ in how they carry context
 - IP context is full header (5 or 6 fields)
 - $< 2^{104}$ permutations, multi-dimensional
 - MPLS context is carried in Label and CoS (Class of Service)
 - $< 2^{20}$ ($*2^3$) permutations, index lookup
 - ATM context is carried in VPI/VCI
 - $< 2^{20}$ permutations, index lookup
 - DiffServ context is carried in "DS" byte
 - $< 2^6$ permutations, direct/index lookup



Bandwidth isolation at Link level

□ ATM

- Non Broadcast Multiple Access (NBMA)
- Awkward match to IP multicast service model
 - LANE BUS
 - RFC 2022 (MARS)
 - other....

□ MultiProtocol Label Switching (MPLS)

- L2 Switching, with L3 *routing controlled* paths
- High efficiency POS links, but ‘router’ QoS needed at merge points
- Re-use of ATM links, but routers still exist at ingress and egress to other packet links

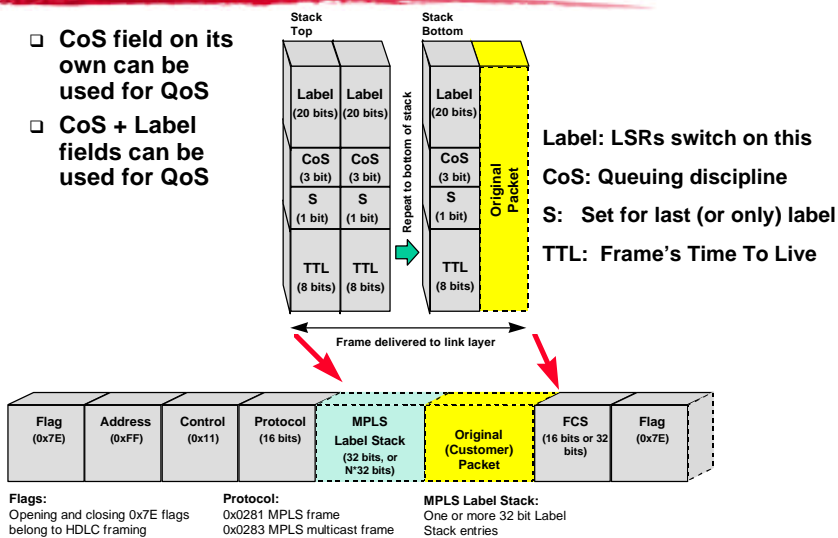
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Inband QoS fields for MPLS over POS

- CoS field on its own can be used for QoS
- CoS + Label fields can be used for QoS



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Bandwidth Isolation..... routers at line rate?



- Various schemes, differing granularity
- DiffServ Routers
 - Per-ToS/DS classification/scheduling behaviors
 - Ingress mapping of IP multicast/tunnel traffic to distinct DS bytes
- Flow isolating Routers
 - Class D traffic can be isolated on a per-group basis (native IP multicast)
 - Per-tunnel isolation (classify on IP Src/Dst)
 - With per-tunnel/group WFQ
 - packet loss, jitter, delay are more controllable
- MPLS Flow Isolation
 - LSPs can be used as tunnel substitutes

Difficulties and Trade Offs



- Classification
 - High speeds traditionally require simple mechanism
- Queuing
 - Number of queues and their management is a problem at gigabit rate
- Scheduling
 - Required algorithms (e.g. Weighted Fair Queuing, WFQ) only recently implementable in hardware for variable packet sizes

DiffServ Compromise

Edge: IP header classification
Core: DS-byte classification
Queues: Limited in the core
Traffic Engineering: Unspecified

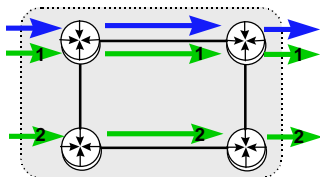
MPLS Compromise

Edge: IP header classification
Core: Label [+CoS] classification
Queues: Unspecified
Traffic Engineering: Supported

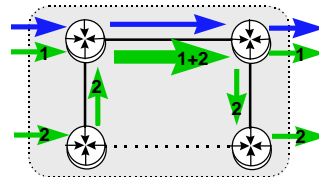
Weakness of limited queues and contexts



- With only limited 'contexts' for customer traffic:
 - Link utilization is engineered for 'normal' operation
- Failure of Edge shaping or WAN connectivity can re-route traffic in unpleasant ways
 - During the re-route, remaining path sees unicast traffic and multicast tunnels sharing queues in excess of engineered limits
 - Service deteriorates for all traffic on same "level"



Flow 1 and 2 have equal priority level - flow 2 is a multicast tunnel



When re-routed, Flows 1 and 2 map to the same queues

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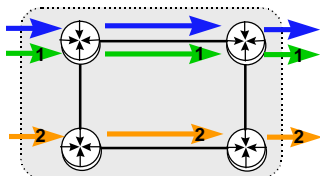
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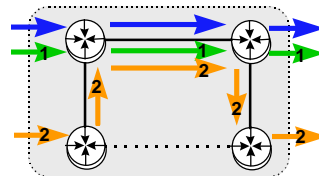
Network Robustness: More Queues and Contexts



- If tunnels are mapped onto distinct 'contexts'
 - During the re-route, remaining path recognizes distinct 'context' for each tunnel and assigns distinct queues
 - *Potential service deterioration can be 'quarantined' to affect only the re-routed traffic*
- Line-rate WFQ across many queues is key to robust QoS architectures



Flows 1 and 2 have different contexts (e.g. full classification of encapsulating IP header)



When re-routed, Flows 1 and 2 map to different queues. Scheduling ensures isolation

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Conclusion



- **Will DiffServ and MPLS help IP Multicast?**
 - YES, although greater granularity preferable

- **Are gigabit routers developing per-tunnel and per-LSP QoS capabilities?**
 - YES

- **Is there a good future for IP Multicast QoS?**
 - YES